

**Specifying Downloadable  
Properties for Reusing Software  
Components: A Case Study of Java**

*Haruhiko Kaiya & Kenji Kaijiri*

Shinshu University, JAPAN

Sep. 13, 2000

# Outline

- The **Role of Specification** for Components
- **Downloadable Components** in Java1
- **Properties** of Downloadable Components
  - code deployment, class load, security manage.
- **How to specify** Downloadable Components
- The **Advantages** of our Specification
  - An example: RMI with Cracking Code
- Conclusion and Future Works

# Role of Spec. for Components

- Spec. as **Manual** for suitable reuse
- Understanding its functionality and limitations
- Base for component browser in IDE
  
- Natural Languages: long, vague.
- Formal Notation
  - Compact, consistent, formal reasoning
  - Hard to write
  - Retrieve the cost by using it repeatedly

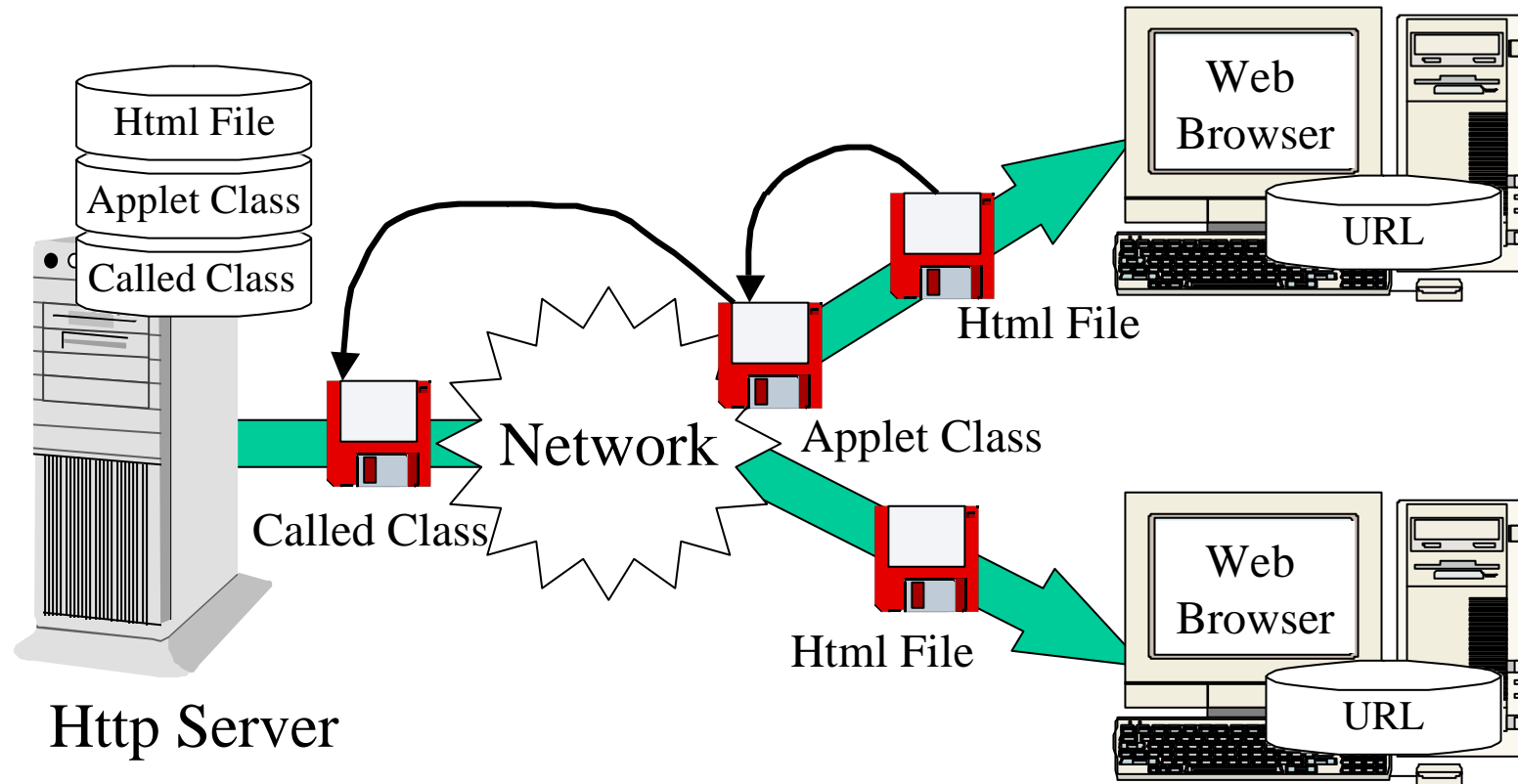
# Traditional Functional Spec.

- **Signatures:** type of arguments, returns and method name – No semantics!
- **Pre and Post Condition** for each method
- **Invariant** for a class
- Enough to specify normal components.
- Not enough for downloadable components

# Downloadable Components

- Loaded from *remote machines*
- *Dynamically* loaded and linked.
- Examples: Java Applets, RMI stub and skeleton.
  
- *Not fully trusted*, i.e. in the *sandbox*.
- Depending on the **services and environments** *outside* the system.

# An example: Applet



# Func. Spec. for D.L. Components

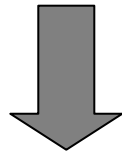
- Traditional Spec: signatures, pre/post conditions, invariants.

+

- **Deployment of bytecodes** over the network
- **Class loader's Policy**: search path for bytecodes.
- **Security Manager** for a Machine: checklist for accessing system resources.

# The Story of this Example

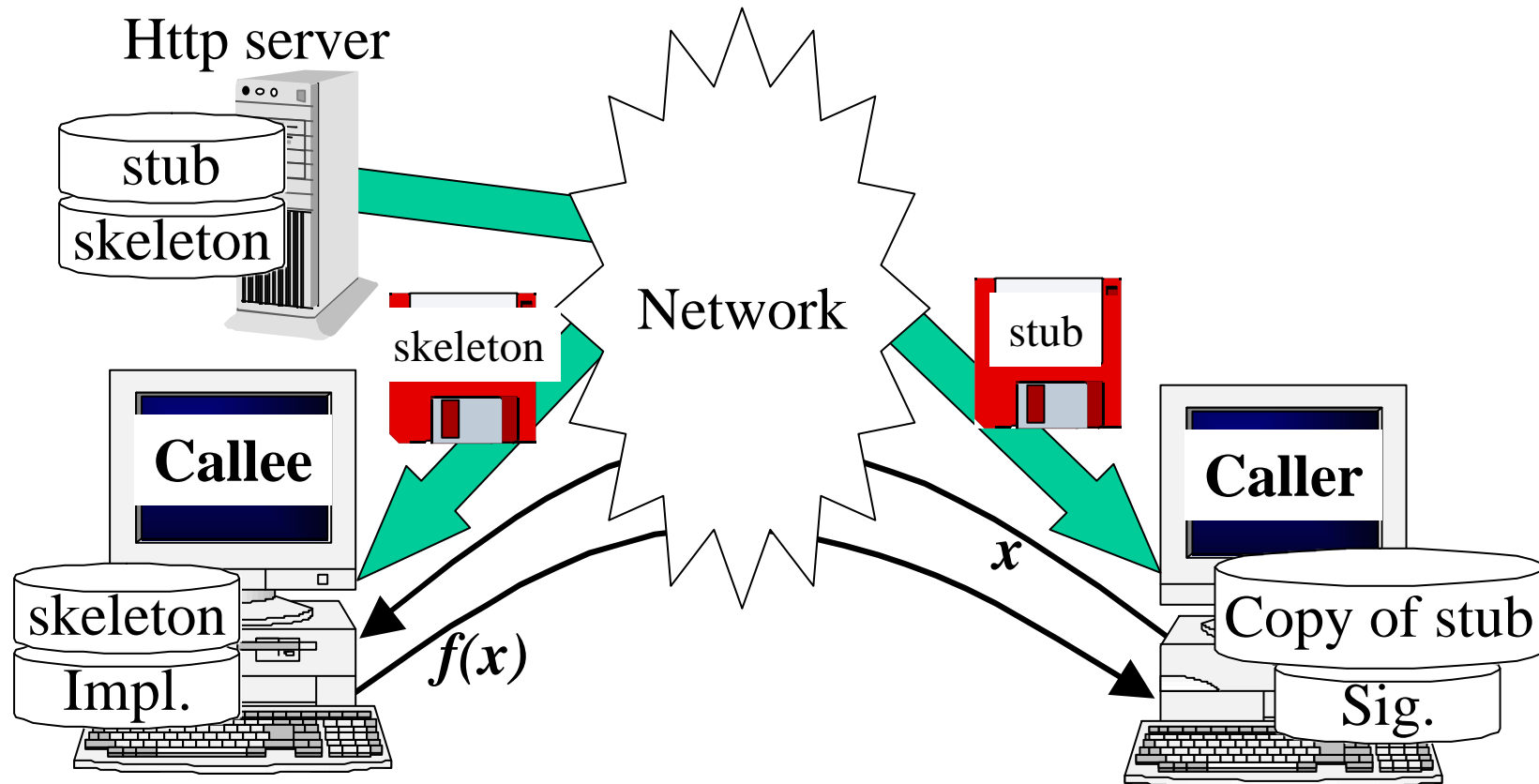
- RMI *with cracking code* in its stub!
- We know, but we *should use* it.
- Risk for *the progress of cracking*.



- Set Security Manager.
- Deploying the copy of current stub in the local.



# Overview of this Example



# Insufficient Specification 1

- Only account for security manager.

*SysRes*

$res : R \leftrightarrow Bool; limit : \mathbb{P} R$

$limit \subseteq \text{dom } res; \forall x : limit \bullet (x, false) \in res$

*SetLimit*

$\Delta SysRes; l? : \mathbb{P} R$

$limit \neq \emptyset \Rightarrow l? = limit'$

# Insufficient Specification 2

- Method body and Cracking Code

$$\frac{\text{Func}}{\frac{x?, y! : \mathbb{Z}}{y! = f(x?)}}$$

$$\frac{\text{Crack}}{\frac{pas! : R; \exists SysRes}{(pas!, true) \in res}}$$

$$F \equiv (Crack \circ Func \wedge \exists SysRes) \setminus \{pas!\}$$

# Formal Reasoning

- This schema tells '*Cracking is established even if the security manager is set*'.

*SetLimt*  $\circ$  *Crack*  $\circ$  (*Func*  $\wedge$   $\exists$ *SysRes*) | *pas!*  $\in$  *l*?

- This is inconsistent.
  - Both (*pas!*, true) *in res* and (*pas!*, false) *in res*
  - *res* is function

# Discussion for insufficient spec.

- In fact in this stage, *Cracking is established!*
- *Security manager is helpless* because of the copy of *stub* **in the local** system.
- *Deployment of bytecodes* is ignored.
- *Class loader's policy* is also ignored.

# Sufficient Specification 1

- Deployment of bytecodes

|  $deploy : Loc \leftrightarrow \mathbb{P} \text{ByteCode}$

- *SysRes* with current location.

*SysRes*

$res : R \leftrightarrow \text{Bool}; limit : \mathbb{P} R; here : Loc$

$limit \subseteq \text{dom } res$

$\forall x : limit \bullet (x, \text{false}) \in res$

$here \in \text{dom } deploy$

# Sufficient Specification 2

- Rel. between the component and its src

<i>Class</i>
$birth : Loc; byte : ByteCode$ $lslctr : seq Loc$
$birth \in \text{ran } lslctr$ $birth \in \text{dom } deploy$

<i>SetLoader</i>
$sl?; seq Loc; \Delta Class$
$lslctr' = sl?$ $\forall x, y : \mathbb{N} \bullet byte \in deploy lslctr' x \wedge$ $x \in \text{dom } lslctr' \wedge lslctr' y = birth' \Rightarrow y \leq x$

# Sufficient Specification 3

- Spec. of Cracking code.

$$\textit{Crack} \equiv [pas! : R; \exists SysRes; \exists Class | \\ \textit{here} \neq \textit{birth} \Rightarrow (pas!, rmtrue) \in res]$$



# Formal Reasoning

- This schema also tells *'Cracking is established even if the security manager is set'*

$$\begin{aligned} & \text{SetLimit } \circlearrowleft (\text{SetLoader} \wedge \exists \text{SysRes } \circlearrowleft \text{Crack } \circlearrowleft \\ & \quad \text{Func} \wedge \exists \text{Class} \wedge \exists \text{SysRes}) \setminus \text{Class} \\ & | \text{pas!} \in l? \wedge sl? = \langle \text{here}, \text{there} \rangle \end{aligned}$$

- The above becomes consistent under the deployment:

$$\text{deploy} = \{(\text{here}, \{\text{byte}, \dots\}), (\text{there}, \{\text{byte}, \dots\}) \dots\}.$$

# Conclusion

- We extend traditional functional spec. for fitting downloadable components in Java1.
- Additional Concept:
  - Security Policy and Management.
  - Class loading
  - Deployment of bytecodes.
- Suitable reasoning for component use.

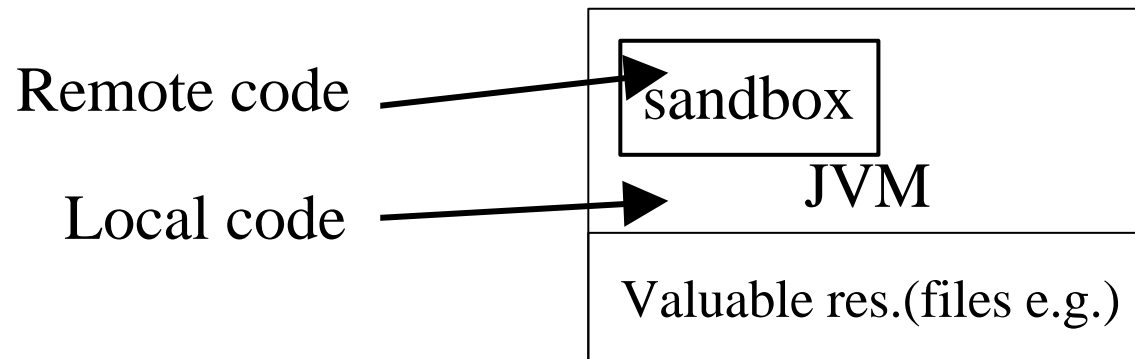
# Future Works

- Extending our Spec. for Java2
  - Code signing (since JDK1.1)
  - Permission and AccessController Class
  - Extended sandbox model
- Other Security Techniques: Proof Carrying...
- Spec. for Scripting, VBscript, JavaScript ....
- Prog. Lang. Independent Spec.
- Embedding our spec. into component browsers

# *Appendix*

# Sandbox Security Model (Java1)

- *Local code* is trusted to have *full access* to vital system resources
- While *downloaded remote code* is not trusted and can *access only the limited resources* provided inside the *sandbox*



# Class Loader (Java1)

- System for loading ByteCodes into JVM
- System Class Loader:
  - From local file system
  - Always trusted.
- User Defined Class Loader:
  - From any sources, e.g. remote systems, byte streams, databases.
  - Normally untrusted.
  - Customizable by programmers.

# Security Manager

- Check List for accessing system vital resources.
  - Read/write files system, create net. connection, create other class loaders.....
- Implemented as a class.
- Methods for accessing res. in API refer the lists.